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# Quantum Information and Quantum Computer

SS 2018

7. Exercise

07.06.2018

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## Task 1

- Use the components of the discrete universal quantum gate (Hadamard, phase, “ $\pi/8$ ” and CNOT gates) to design a circuit for the quantum Fourier transform of a 3-qubit register.
- What is the explicit form of the transformation matrix for the circuit in a)?

## Task 2

The *order* of an integer number  $y$  modulo a (natural) number  $N > y$  is defined as the smallest positive integer power  $r$  fulfilling

$$y^r = 1(\text{mod } N) \quad .$$

An important theorem says that for a number  $N$  which can be factorised it is *very likely* that the order  $r$  is even and  $y^{r/2} \neq \pm 1(\text{mod } N)$  is valid, if the largest common factor of  $y$  and  $N$  is equal to 1 ( $\text{GGT}(y, N) = 1$ ). In this case one finds due to

$$\begin{aligned} y^r &= (y^{r/2})^2 = 1(\text{mod } N) \\ (y^{r/2})^2 - 1 &= 0(\text{mod } N) \end{aligned}$$

the relation  $(y^{r/2})^2 - 1 = jN$  with  $j = 1, 2, 3, \dots$ . On the other hand, the relation  $(y^{r/2})^2 - 1 = (y^{r/2} + 1)(y^{r/2} - 1)$  leads to  $\text{GGT}(y^{r/2} + 1, N) \cdot \text{GGT}(y^{r/2} - 1, N) = N$ . In other words,  $\text{GGT}(y^{r/2} + 1, N)$  and  $\text{GGT}(y^{r/2} - 1, N)$  are the factors of  $N$ . Check this claim based on the concrete example  $y = 10$  and  $N = 21$ , i. e. determine the order of  $y$  modulo  $N$  and with its help the factors in 21.

## Task 3

Using the Shor algorithm for factorizing the number  $N = 21$  with the (arbitrary) choice  $y = 11$  and adopting 9 qubits yielded in one run on a quantum computer the result 427. Check whether the correct factors  $p$  and  $q$  can be obtained from this result.

(Please, turn the page ...)

**Task 4**

In order to apply the phase-estimation algorithm it is necessary to prepare the corresponding eigenvectors. In the context of the Shor algorithm these are the eigenvectors of the operator  $\hat{U}_x$  that fulfils

$$\hat{U}_x |y\rangle = |xy \pmod{N}\rangle$$

with  $y \in \{0, 1\}^m$ . Here, the additional convention  $xy \pmod{N} \equiv y$  for  $N \leq y \leq r - 1$  (where  $r = 2^m$ ) was adopted.

- a) Demonstrate that for integer numbers  $0 \leq s \leq r - 1$

$$|u_s\rangle = \frac{1}{\sqrt{r}} \sum_{k=0}^{r-1} e^{-2\pi i sk/r} |x^k \pmod{N}\rangle$$

are the eigenvectors of the operator  $\hat{U}_x$ , if the order (see previous task) of  $x \pmod{N}$  is equal to  $r$ .

- b) The efficient preparation of the eigenvectors given in a) appears on the first glance impossible. Instead of a preparation of the single eigenvectors  $|u_s\rangle$  the relation

$$\frac{1}{\sqrt{r}} \sum_{s=0}^{r-1} |u_s\rangle = |1\rangle$$

is used. Proof the validity of this relation.

- c) Discuss the consequences of the trick for avoiding the preparation of single eigenvectors described that was described in b) has for the Shor algorithm.